

Common Points on Elliptic Curves: The Achilles' Heel of Fault Attack Countermeasures

Alberto Battistello

April 12, 2014

- 1 Introduction
- 2 A Simple Fault Attack
- 3 Our New Attack
- 4 Conclusion

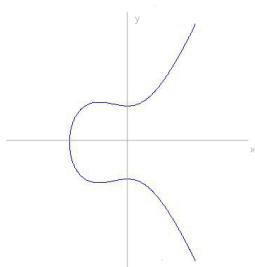
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Short Weierstrass Equation

Curve Equation

$$\mathcal{E}(\mathbb{F}_p) : y^2 = x^3 + \mathbf{a} \cdot x + \mathbf{b} \bmod p$$

$$P = (x_p, y_p)$$



Curve Equation

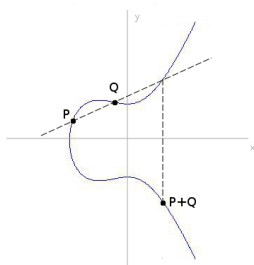
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Addition/Doubling

- $P + Q = f_A(P, Q, \mathcal{E})$
- $P + P = f_D(P, \mathcal{E})$

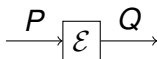
b is never used.



$$\mathcal{E}(\mathbb{F}_p) : y^2 = x^3 + a \cdot x + b \pmod{p}$$

Scalar Multiplication

$$[d]P = \underbrace{P + P + P + P}_{d \text{ times}} = Q$$



ECDLP

Given $P, Q \in \mathcal{E}(\mathbb{F}_p)$ find d such that $Q = [d]P$.

The ECDLP complexity depends on the order of P over $\mathcal{E}(\mathbb{F}_p)$.

Order of Point

$$\text{Ord}_{\mathcal{E}}(P) = \prod q_i^{\alpha_i}.$$

$\max q_i \leq 112\text{-bit} \Rightarrow \rho\text{-Pollard.}$

The same point on another curve may have a weaker order.

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A Simple Fault Attack

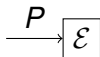
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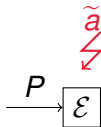
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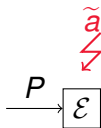
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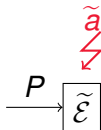
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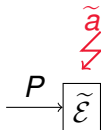
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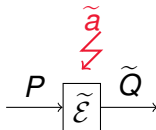
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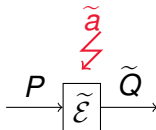
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- 6 $\tilde{Q} \Rightarrow d$



Two simple countermeasures

Parameters checksums

Checksum of all parameters (a, b, p, x_p, y_p, \dots) checked before output.

$$CRC(\tilde{a}) \neq CRC(a)$$

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Point on curve test

- Initial test $y_p^2 = x_p^3 + a \cdot x_p + b \pmod p$
- Final test $y_q^2 = x_q^3 + a \cdot x_q + b \pmod p$

$$y_q^2 \neq x_q^3 + \tilde{a} \cdot x_q + b \pmod p$$

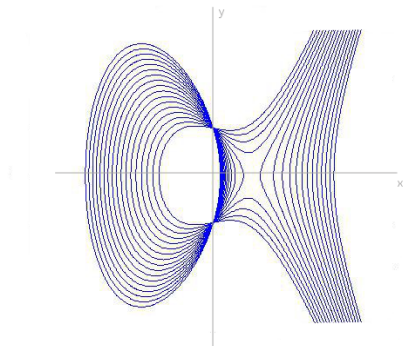
The two tests have been considered equivalent . . .
up to NOW

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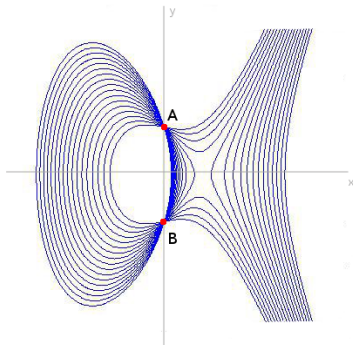
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Common Points:

$$A, B : (0, \pm\sqrt{b})$$



Attack PoC tests

$$\mathcal{E}(\mathbb{F}_p) : y^2 = x^3 + a \cdot x + b \bmod p$$



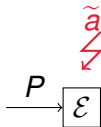
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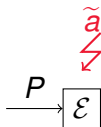
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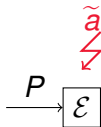
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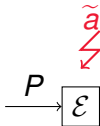
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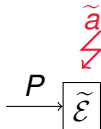
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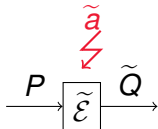
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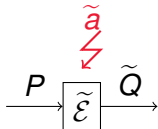
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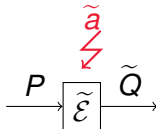
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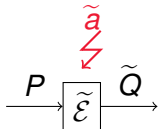
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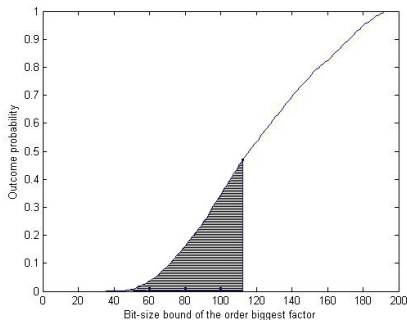


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- 6 $Ord_{\tilde{\mathcal{E}}}(P) \neq Ord_{\mathcal{E}}(P)$
- 7 $\tilde{Q} = [d]P \leftarrow$ **ECDL**



Faults producing a weak curve.



Probability that the biggest factor is smaller than ECDLP record of 112 bit is more than 45%.

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Common Points

Common Points have never been remarked before.

PoC vs Checksums

The two countermeasures are NOT equivalent!

Curves

Parameter b must be a square.

- NIST: P-192, P-256, P-384.
- Secp: 192r1, 256r1, 384r1, 521r1.

Thank you
Feel free to ask any question

